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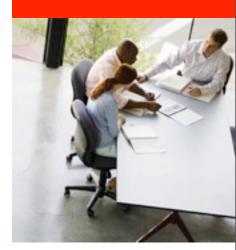
Using DTrace for Leverage

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The many faces of DTrace

- Solve customer escalations
- Debugging
 - Before integration
 - After integration
- Gross performance analysis

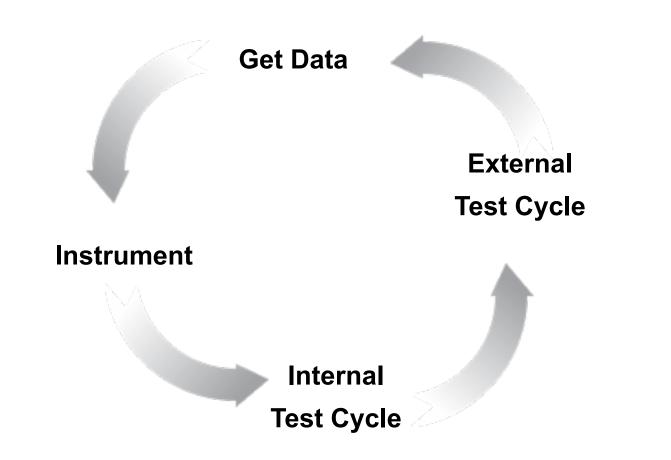




Section Customer Escalations



Getting debug code to customers



Solving Customer Escalations

Gather data at a customer site

- Raw analysis
- Predictive confirmation
- Reproduce problem in-house
 - How do we know we have captured everything?
 - Can we see what the customer is seeing?
- Get a fix in place
 - How do we know the fix works?
 - How do we convince the customer?
- Provide the fix to the customer
 - Predictive confirmation
 - Emergency backup what data to gather?

Preventing Customer Escalations

- Do not want to provide a kernel that causes problems
- Do not want to provide a kernel that stays in production
 - You fixed a specific issue
 - Customer wants to stay on that "release"
 - Your support department may not know how to handle that release



Preventing Customer Escalations

- Do want to predict what the customer will see
- Do want to confirm the fix
- Do want to get them on the path to a supported release



6882460 - mountd storms

- Customer has observed that simultaneously mounting shares via NFSv3 and immediately beginning to read data--with hundreds to near 1,000 Linux clients-causes some mounts to fail.
- NFSv4 does not exhibit the same issue
- waiting for a short time between mounting and reading also alleviates the errors.
- <u>http://bugs.opensolaris.org/bugdatabase/</u> view_bug.do?bug_id=6882460

Usual suspects

- Packet traces show either
 - Quick responses to mount requests
 - No response at all
- Name server is quick to respond
 - Putting names in /etc/hosts does not improve things
- No netgroups
- Bumping number of mountd threads improves things

Recreating a mountd storm

- Want to be able to blast N mount requests to a server
- Do not care about the replies
- Need to have both forward and reverse DNS entries
- Need to consider one request per IP versus many
 - Only first request should see a name lookup across the wire
 - Case appears to be insensitive to name lookups
- Use a Perl script to blast N UDP mount requests
 - Are serialized, but end up arriving quick enough
 - Easy enough to scale past customer needs
 - Might need to really scale in order to get meaningful data

Time to drill down

```
#!/usr/sbin/dtrace -Fs
```

```
/*
 *
        # ./mountd.d `pgrep -x mountd`
 */
dtrace:::BEGIN
ł
       printf("Sampling... Hit Ctrl-C to end.\n");
}
pid$1::mount:entry
{
       self->timestamp[probefunc] = timestamp;
       @function count[probefunc] = count();
       self->trace = 1:
}
pid$1::mount:return
ł
       @function quantize[probefunc] =
             quantize(timestamp - self->timestamp[probefunc]);
       self->timestamp[probefunc] = 0;
}
```

Real customer data

mount		
449		
mount		
value	Distribution	count
2097152		0
4194304	0@	26
8388608	0	14
16777216	0	13
33554432	0	9
67108864	0@	18
134217728	0@	28
268435456	@@@@@	59
536870912	0@@	31
1073741824		5
2147483648		4
4294967296		2
8589934592		2
17179869184	0@@	30
34359738368	@@@@@@@@@@@@@@@@@@	208
68719476736		Θ



- 1000 mount requests
- 449 made it to mount(), which means 551 did not
- 30 mount requests took between 17-34 seconds
- 208 mount requests took between 34-68 seconds



Reproduction with 4000 requests

washdc# ./mountd.d `pgrep -x mountd` dtrace: script './mountd.d' matched 3 probes CPU FUNCTION Sampling... Hit Ctrl-C 5 | :BEGIN to end. ^C mount mount ----- Distribution ------ count value 0@@@@@@@@@@@@ @ 268435456 |@@@@@@@@

Comparing real vs synthetic

- Went to 4000 to get more data
 - Went up to 40k at some points
- Lab machine appears faster
 - Did not model NFS traffic, just mountd
 - Did do a NFS load later
- Mostly happy we simulated a gap
 - I.e., data appears representative of the general problem

Looking at the mount() code

- caller
 - Picked apart request
 - Determined which procedure to invoke
- do name lookup
- do access check
- reply to client
- log into BSM audit trail
- add to /etc/rmtab
- return

Auditing looked suspicious

- Fast replies once we got them
- Increasing threads appeared to alleviate problem
 - Actually provided a queue for starvation
- We were starved waiting on a global lock
- Applied a fix to add entries to BSM audit trail asynchronously

Asynch auditing with 4000 requests

washdc# ./mountd.d `pgrep -x mountd` dtrace: script './mountd.d' matched 5 probes CPU FUNCTION Sampling... Hit Ctrl-C 0 | :BEGIN to end. ^C 2970 mount mount value ----- Distribution ----- count 32768 0 2051 131072 | @@@@@@@@@@@@@@ 902 262144 15 524288 1 1048576 1 2097152 0

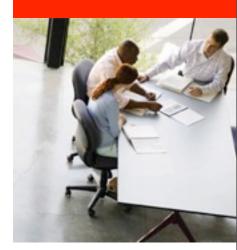
Analysis of async auditing

- Service about 3 times as many mount requests
- Long pole is about 130 microseconds
 - versus 0.5 seconds
- We were also able to show it was independent of the number of mountd threads
 - Get the customer off of this customization!

Customer results with async logging

mount

value	Distribution	count
16384		0
32768		4
65536	00000	123
131072	@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@	530
262144	@@@@@@@@@@	232
524288	0@@	66
1048576		7
2097152		7
4194304		0
8388608		3
L6777216		0



Section Debugging



Debugging Before integration

- Live system debugging
 - Save mdb for core dumps
- Want to avoid printf()
- Want to avoid recompilation
- What can existing function calls tell you?
- Where do you need to add a static probe?

Adding static probes

- Counters
- Error detection
 - Avoid logging to console
 - Certainly avoid spamming console/logs
- Allow targeted data gathering
 - Avoid ifdef
 - Avoid /etc/system
 - Allow dtrace to be your on switch

Debugging After integration

- Help others understand your code
 - Provide scripts to allow them to debug problems
- Help you understand if more static probes needed
 - First external use
 - Customers will have same needs



Debugging spe

- spe assigns layouts
 - mds does not need to use it
 - Loads from
 - /etc/npools.spe
 - /etc/policies.spe
- Always being asked
 - Did spe work?
 - nfsstat -I can tell you that
 - How do I know which policy fired?
- A small script can tell all

spe.d part 1

```
nfsv4:::spe-i-check open
{
        printf("%d (%d) from %s is checking %s",
            (uid t)arg0, (gid t)arg1, stringof(arg3), stringof(arg2));
}
nfsv4:::spe-i-policy eval
{
        printf("Policy %d %s with error %d from %s",
            (uint t)arg0, (boolean t)arg1 ? "matched" : "did not
match",
            (int)arg2, stringof(arg3));
}
::nfs41 spe allocate:entry
        self->addr = (struct netbuf *)arg2;
        self->stripe count = (count4 *)arg4;
        self->unit size = (uint32_t *)arg5;
        self->mds sids = (mds sid **)arg6;
        self->loaded sids = 0;
}
```

spe.d part 2

```
::nfs41 spe allocate:return
/args[1] == 0/
        printf("Policy has %d stripes and %u block size",
            *self->stripe count, *self->unit size);
::nfs41 spe allocate:return
/args[1] != 0/
ł
        printf("No matching policy");
}
::mds_ds_path_to_mds_sid:entry
        self->ustring = (utf8string *)arg0;
        self->ss name = stringof(self->ustring->utf8string val);
        self->mds sid = (struct mds sid *)arg1;
}
```

spe.d part 3

```
::mds ds path to mds sid:return
/args[1] == 0/
        ss name = (char *)alloca(self->ustring->utf8string len + 1);
        bcopy(self->ustring->utf8string val, ss name,
            self->ustring->utf8string len);
        ss name[self->ustring->utf8string len + 1] = ' 0';
        printf("mds sid[%d] = %s", self->loaded sids++,
            stringof(ss name));
}
::mds_ds_path_to_mds_sid:return
/args[1] != 0/
Ł
        ss name = (char *)alloca(self->ustring->utf8string len + 1);
        bcopy(self->ustring->utf8string val, ss name,
            self->ustring->utf8string len);
        ss name[self->ustring->utf8string_len + 1] = '\0';
        printf("ERROR - could not find a matching pgi for %s",
            stringof(ss name));
}
```

Policies not loaded

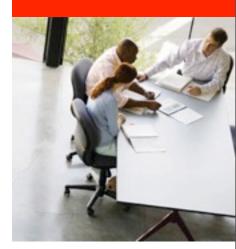
[root@pnfs-minipit2-5 ~]> ./spe.d dtrace: script './spe.d' matched 8 probes CPU ID FUNCTION:NAME 1 28427 nfs41_spe_allocate:spe-i-check_open 60001 (60001) from 10.1.235.62 is checking /diskpool/DS/P2/foo 1 59258 nfs41_spe_allocate:return No matching policy

DS dataset does not match mds

1 3420 nfs41 spe allocate:spe-i-check open 200096 (10) from 10.1.233.191 is checking /diskpool/JUNK/TEST/P5/tomper 1 3419 nfs41 spe allocate:spe-i-policy eval Policy 101 did not match with error 0 from 10.1.233.191 1 3419 nfs41 spe allocate:spe-i-policy eval Policy 102 did not match with error 0 from 10.1.233.191 1 3419 nfs41 spe allocate:spe-i-policy eval Policy 103 did not match with error 0 from 10.1.233.191 1 3419 nfs41 spe allocate:spe-i-policy eval Policy 104 did not match with error 0 from 10.1.233.191 1 3419 nfs41 spe allocate:spe-i-policy eval Policy 111 matched with error 0 from 10.1.233.191 1 63756 mds_ds_path to mds sid:return mds sid[0] = pnfs-minipit1-6:pNFSpool1/p1DS2 1 63756 mds ds path to mds sid:return mds sid[1] = pnfs-minipit1-6:pNFSpool2/p2DS2 1 63756 mds_ds_path_to_mds_sid:return mds_sid[2] = pnfs-minipit1-6:pNFSpool3/p3DS2 1 63756 mds ds path to mds sid:return mds sid[3] = pnfs-minipit1-7:pNFSpool1/p1DS1 1 63756 mds ds path to mds sid:return 1 57043 nfs41 spe allocate:return No matching policy

What went wrong?

- Policy stated that 5 datasets were needed
 - 111, 5, 32k, default, path == /diskpool/JUNK/TEST/P5
- Only 4 found
- DS has a dataset named:
 - pnfs-minipit1-7:pNFSpool2/p2DS1
- /etc/npools.spe has dataset named as
 - pnfs-minipit1-7:pNFSpool1/p2DS1



Section Gross Performance



Gross performance analysis

What is the performance cost of a

- New feature
- Bug fix
 - What was the performance before?
- "Gross" means
 - A metric you can use as a developer
 - Gives you a rule of thumb approximation
 - Not something you might put in a formal report
- But you need to be able to answer the question about the impact

Gross Performance of referrals

- Referrals built on top of mirrormount framework
- As part of integration, asked ourselves
 - What is the performance of doing a referral?
- Note: we never asked ourselves that when we did the mirrormount work

Timing referrals

- Kernel versus userland
- But pretty much the same as before

First look at the performance

4098						
nfs4_trigger_mount						
value Distribution cour	١t					
4096 0						
8192 19						
16384 @@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@@)					
32768 27						
65536 2						
131072 0						
262144 1						
524288 1 1						
1048576 0						
2097152 0						
4194304 0						
8388608 @@@@@@@ 672						
	5					
33554432 @ 95						
67108864 [@ 102						
134217728 @ 97						
268435456 @ 74						
536870912 42						
1073741824 1						
2147483648 0						

Okay, 1/2 second sounds gross

- About 750 referrals, all under one parent
- mirrormounts gave about the same numbers
- We did get a 68 second outlier



3099 mirrormounts under 1 real

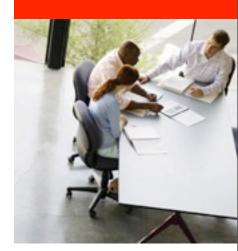
3099	igger_mount		
nfs4_trigger	_mount		
value		Distribution	count
1048576			0
2097152			1
4194304			37
8388608	0@@		235
16777216	0@		147
33554432	0@		189
67108864	0@		150
134217728	0@@		211
268435456	0@@		234
536870912	0@@		224
1073741824	0@@@		329
2147483648	0@@@@		420
4294967296	0@@@@		375
8589934592	0@@		265
17179869184	0@@		198
34359738368	0		84
68719476736			0

2049 mirrormounts under 1051 real

nfs4_t	rigger_mount			
2049				
nfs4_trigger_mount				
value	Distribution	count		
8388608		0		
16777216		9		
33554432	000	132		
67108864	@@@@@@@@@@@@@@@	789		
134217728	@@@@@@@@@@@@@@@@@	874		
268435456	0000	211		
536870912	@	33		

Analysis - Another locking issue

- We are holding a lock as we do a syscall to perform an ephemeral mount
 - I.e., mirrormount or referral
- Lock prevents tree from going away from other events
- With 1051 real mount points, we distribute the load
- With 1 real mount point, we see the strain
- Can model a reader/writer lock to alleviate the starvation



Section Take home message



Take it home with you

- DTrace allows you to avoid printf()
 - I.e., no recompiling
- DTrace allows you to gather data at a customer's site
 - They can even gather it. :->
- No instrumented kernels polluting a customer's site

(i) ORACLE IS THE INFORMATION COMPANY