

NFSv4 Performance Across A WAN

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Part I: State of the WAN Performance Art

- Our assumptions
- Benchmarks
- Part II: Strategies for improving WAN Performance
 - Helping all versions of NFS
 - What's in NFSv4 today
 - The Future
- Questions



• Why is NFS on a WAN so slow?

"I'm a telecommuter. I find NFSv3 performance across the country to suck."

"Whenever I use NFS from home, all other network activity slows to a crawl."

"Why can't I quickly access my colleagues' files on a remote file server?"



Remote shared intermediate cache

- Doesn't help files accessed by only a single client (home directory)
- Doesn't help single remote clients
- Doesn't help read-once workloads
- Still requires GETATTRs and LOOKUPs to maintain close-to-open cache consistency



• WAN performance is slow because:

- The client waits for attribute cache validation
- The client waits for access authorization
- The client has purged its cache and must re-read data or metadata
- The client waits for reads or writes to stream across a slow link
- The client waits for network congestion to clear



NFS congestion test

- Is the NFS client a friendly network neighbor on a slow link?
- Does it hamper its own performance?
- Does delegation help?
 - Application performance comparison
 - Reduction in network traffic?
- Time wait analysis
 - Which RPCs cause the most wait time?



- Using the "sync" mount option on Linux appears to reduce network congestion
 - "Sync" allows only one read or write on the network at a time
- Quantify WAN link congestion due to NFS client
 - Can such congestion be reduced or eliminated by careful client design?
 - Does client detect congestion it causes?





- One DSL link with VPN tunnel
- Four systems, total:
 - Local end: NFS client and ping system
 - Remote end: NFS server and ping target host
- Two traffic streams:
 - NFS writes from client system to server
 - Pings from ping system to ping target
- Quiescent ping RTT = 75-80ms
- TCP window size on uplink: 10KB



Allow one 4K write at a time

- Ping RTT_{avg} = 185ms, RTT_{max} = 384ms
- 1:59 elapsed to write 1363656 bytes
- Allow two 4K writes at a time
 - Ping RTT_{avg} = 396ms, RTT_{max} = 550ms
 - 1:32 elapsed
- Allow three 4K writes at a time
 - Ping RTT_{avg} = 671ms, RTT_{max} = 1004ms
 - 1:32 elapsed



Allow one 8K write at a time

- Ping RTT_{avg} = 315ms, RTT_{max} = 629ms
- 1:44 elapsed to write 1363656 bytes
- Allow two 8K writes at a time
 - Ping RTT_{avg} = 920ms, RTT_{max} = 1105ms
 - 1:32 elapsed
- Allow three 8K writes at a time
 - Ping RTT_{avg} = 1479ms, RTT_{max} = 1661ms
 - 1:32 elapsed



- Smaller r/wsize is more friendly
 - Allows other traffic to interleave
- More than two concurrent write requests doesn't improve throughput
 - Two requests fill TCP window
- Existing congestion control algorithm doesn't help
 - Little RTT variance prevents triggering congestion control



Hypothesis:

- RTT matters for all workloads
- Data throughput matters only for data intensive workloads
- How long does client wait for:
 - Metadata reads (synchronous)
 - Metadata updates (synchronous)
 - Data reads and writes (sync and async)



- Client: Linux 2.4.20 with readdirplus and access (cold cache)
- Server: NetApp F880
- Mount options: vers=3,tcp,rsize=4K,wsize=4K





RPC count (total)





RPC count (meta reads)





- ▶ 84,550 RPCs total
 - Almost half are synchronous metadata reads
- LAN: 84K RPCs * 45% * 1ms = 38 seconds
- WAN: 84K RPCs * 45% * 100ms = 63 minutes



- Prefetch metadata (VFS cooperation)
- More effective cache revalidation
- Parallelize synchronous RPCs (access + getattr, access + lookup)
 - NFSv3: parallel issue, NFSv4: compound RPC
- Depend on "bulk stat"
 - NFSv4: compound RPC
- Delegation



- Solaris NFSv4 with file delegation
- tar cf -/tmp /usr/include
- ~2100 files, ~150 directories



Without delegation

- Cold cache: 10.5 seconds elapsed, ~50K ops
- Warm cache: 5.9 seconds elapsed, ~24K ops
- With delegation
 - Cold cache: 11.2 seconds elapsed, ~50K ops
 - Warm cache: 2.5 seconds elapsed, ~7.7K ops



- After cache is warm, delegation effectively reduces network traffic
- Second run with delegation uses 85% fewer RPCs
- WAN: (cold) 50K RPCs * 100ms = 84 minutes
- WAN: (warm) 7.7K RPCs * 100ms = 13 minutes



- General help for all versions of NFS
- Using new features in NFSv4 (RFC 3010)
- Driving future versions of NFS



Eliminate superfluous network requests

- Reduce GETATTRs and LOOKUPs to bare minimum
- Make good use of READDIRPLUS
- Invalidate data and attribute cache less aggressively
- Reduce round-trips
 - Issue synchronous RPCs in parallel
 - Prefetch attribute data



Use advanced congestion control

- Congestion more broadly defined to include increases in packet latency
- Use TCP window size to control maximum number of concurrent requests



File delegation

- Can greatly reduce READ traffic
- Client can trickle write-backs even after a CLOSE
- Compound RPC
 - More operations per compound means fewer round trips
 - Client VFS architecture limits applicability



Directory delegation

- Purpose: reduce or eliminate LOOKUPs and READDIRs
- Server notifies client when directory entries are created or destroyed



Bulk GETATTR

- Attribute "read ahead"
- Grab attributes for a bunch of files at once
- Much like READDIRPLUS, but not necessarily triggered by getdents()



Questions

