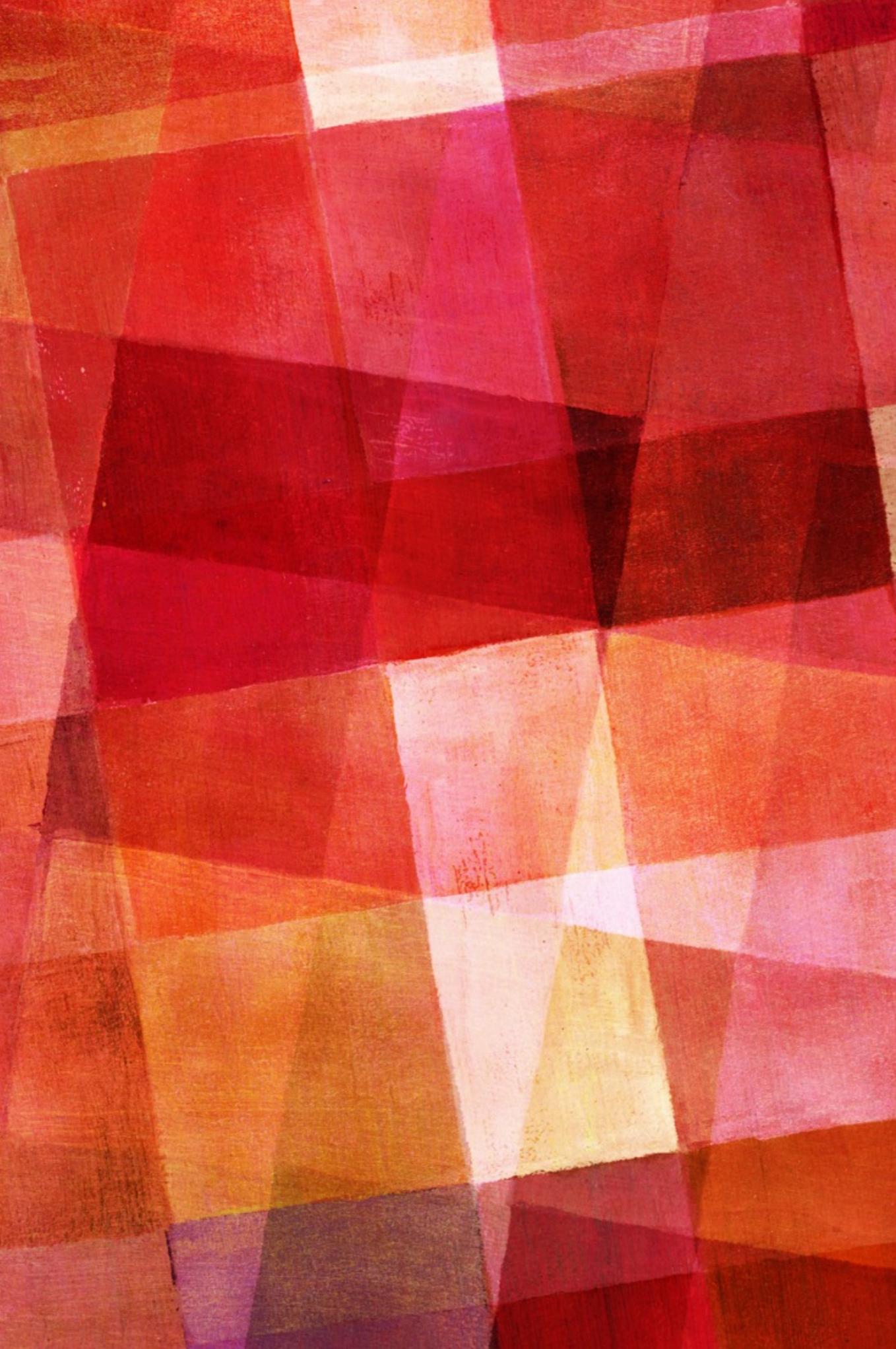


NFS/RDMA BASICS

Part Three – Code Organization





CODE ORGANIZATION

- Client transport overview
- Server transport overview
- NFSv4.1 backchannel operation

RDMA VERBS

- In the Linux kernel, the RDMA verbs API is provided by a set of function calls and data objects.
 - These work with any RDMA-enabled network fabric
- Verb names start with `ib_` :
 - `ib_post_send`, `ib_modify_qp`, `ib_sge`
- RDMA core functionality uses names start with `rdma_` :
 - `rdma_resolve_addr`, `rdma_create_qp`

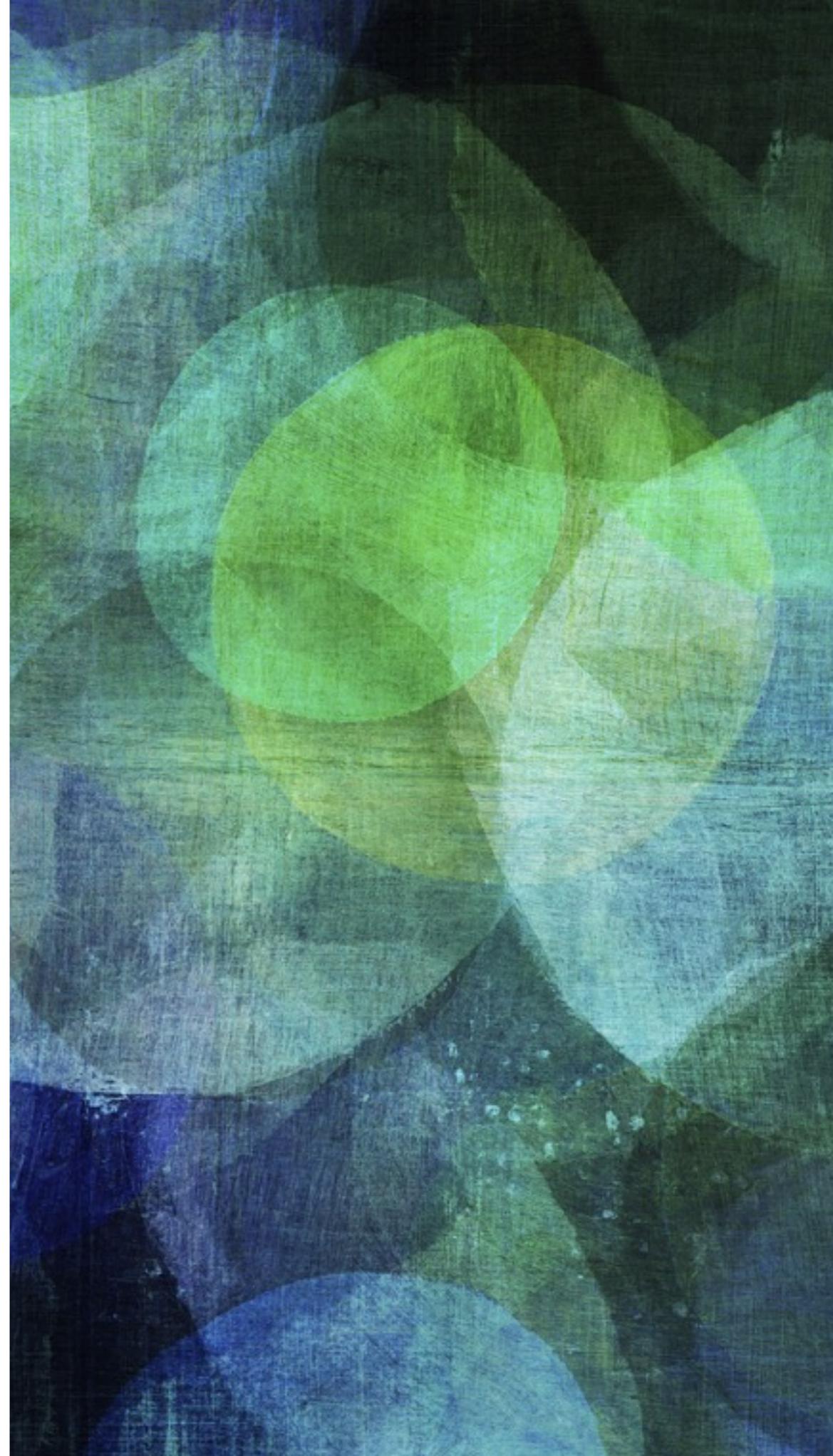
UPCALLS

- In the Linux kernel, verbs provider upcalls are used:
 - When a Send or Receive completion fires
 - When reporting a connection-related event
 - When reporting a QP error
- Upcalls may occur in process or soft IRQ context.
 - Posting a Send or Receive WR can be done in either context.

THE RPCRDMA KERNEL MODULE

- Source code is located in `net/sunrpc/xprtrdma`:
 - Server code is in files names `svc_rdma_*`
 - Client code is everything else
- Currently one module, `rpcrdma.ko`, contains both the client and server transports.

CLIENT TRANSPORT OVERVIEW



CLIENT TRANSPORT SWITCH

➤ transport.c

Method	Purpose
reserve_xprt	Take write lock
release_xprt	Release write lock
connect	Establish a connection
close	Close a connection
buf_alloc	Allocate buffers for Call and Reply
buf_free	Release buffers
send_request	Send an RPC Call
timer	An RPC timeout occurred

MARSHALING RPC CALLS

- `rpc_rdma.c`
 - Main entry point is `rpcrdma_marshal_req`.
 - Decides whether to send each RPC Call inline, use scatter-gather, reduce data items, or use special chunks.
 - Uses NFS XDR reply size information
 - Chunk lists are constructed and memory is registered.
 - The Transport Header is built in a separate buffer, then this buffer plus the buffer containing the RPC message are Sent together.

HANDLING RPC CALLS WITH DATA PAYLOADS

- NFS sets a flag in the `xdr_buf` to indicate when the NFS operation is allowed to use a Read chunk. The `xdr_buf`'s page list contains the data payload.
 - If the RPC Call is smaller than the inline threshold, the data buffer is made part of the Send message, using the Send WR's scatter-gather list.
 - If the RPC Call is large, the data buffer is registered as a Read chunk and added to the Read list.
- If the RPC Call is large and no Read chunk is allowed, the whole message is registered and added to the Read list as a Position Zero Read chunk.

PREPARING FOR RPC REPLIES WITH DATA PAYLOADS

- NFS sets a flag in the `xdr_buf` to indicate when the NFS operation is allowed to use a Write chunk. The `xdr_buf`'s page list contains the data buffer.
 - If the estimated maximum size of the RPC Reply is smaller than the inline threshold, no additional action is needed.
 - If the estimated maximum size of the RPC Reply is large, the data buffer is registered as a Write chunk and added to the Write list.
- If the estimated maximum size of the RPC Reply is large and no Write chunk is allowed, a Reply chunk is registered and added to the Transport Header.

HANDLING REPLIES

- `rpc_rdma.c`
 - Receive upcall runs in soft IRQ context
 - DMA sync, process the credits field, queue work
- `rpcrdma_reply_handler` runs in workqueue context
 - Fully parses the transport header
 - Invalidates and DMA unmaps memory associated with request
 - Pulls up and reconstructs the RPC Reply `xdr_buf`
 - Invokes `xprt_complete_rqst`

GENERIC RDMA HELPERS

- verbs.c
 - Send, Receive, QP error, and connect upcalls
 - Transport set up and tear-down
 - A *regbuf* is a memory region with an lkey and DMA mapping state
 - Registered for local access only
 - Used internally by the transport for RPC buffers
 - Helpers for posting Send and Receive WRs

MEMORY REGISTRATION OPS

- Specific methods for performing memory registration and invalidation on memory that belongs to the upper layer

Method	Purpose
<code>map</code>	Register an MR
<code>unmap_sync</code>	Invalidate all MRs for an RPC
<code>unmap_safe</code>	Invalidate or recover all MRs for an RPC
<code>recover_mr</code>	Recover one MR
<code>open</code>	Compute registration parameters
<code>maxpages</code>	Return maximum pages per MR
<code>init_mr</code>	Prepare one MR for use by the transport
<code>release_mr</code>	Release MR before transport destruction

FRWR MEMORY REGISTRATION

- `frwr_ops.c`
 - Registering memory for one RDMA segment:
 - DMA map the region then post a FastReg WR to register it
 - WR is not signaled
 - Invalidating memory for one RPC:
 - Post LocalInv WRs for all registered MRs
 - Wait for completion
 - DMA unmap all MRs

FMR MEMORY REGISTRATION

- `fmr_ops.c`
 - Registering memory for one RDMA segment:
 - DMA map the region
 - Use `ib_map_phys_mr` to register it
 - Invalidating memory for one RPC:
 - Build a list of all MRs
 - Use `ib_unmap_fmr` to invalidate them
 - DMA unmap all MRs

THE CONNECT WORKER

- transport.c and verbs.c
 - IP address is resolved to a GID/LID (native address)
 - Connecting a transport is serialized with sending RPC Calls
 - Connect worker also handles device unload events
 - Registered memory has to be “re-registered” after a reconnect
 - DMA mapped regbufs have to be remapped after a device unload

INTERESTING DATA STRUCTURES

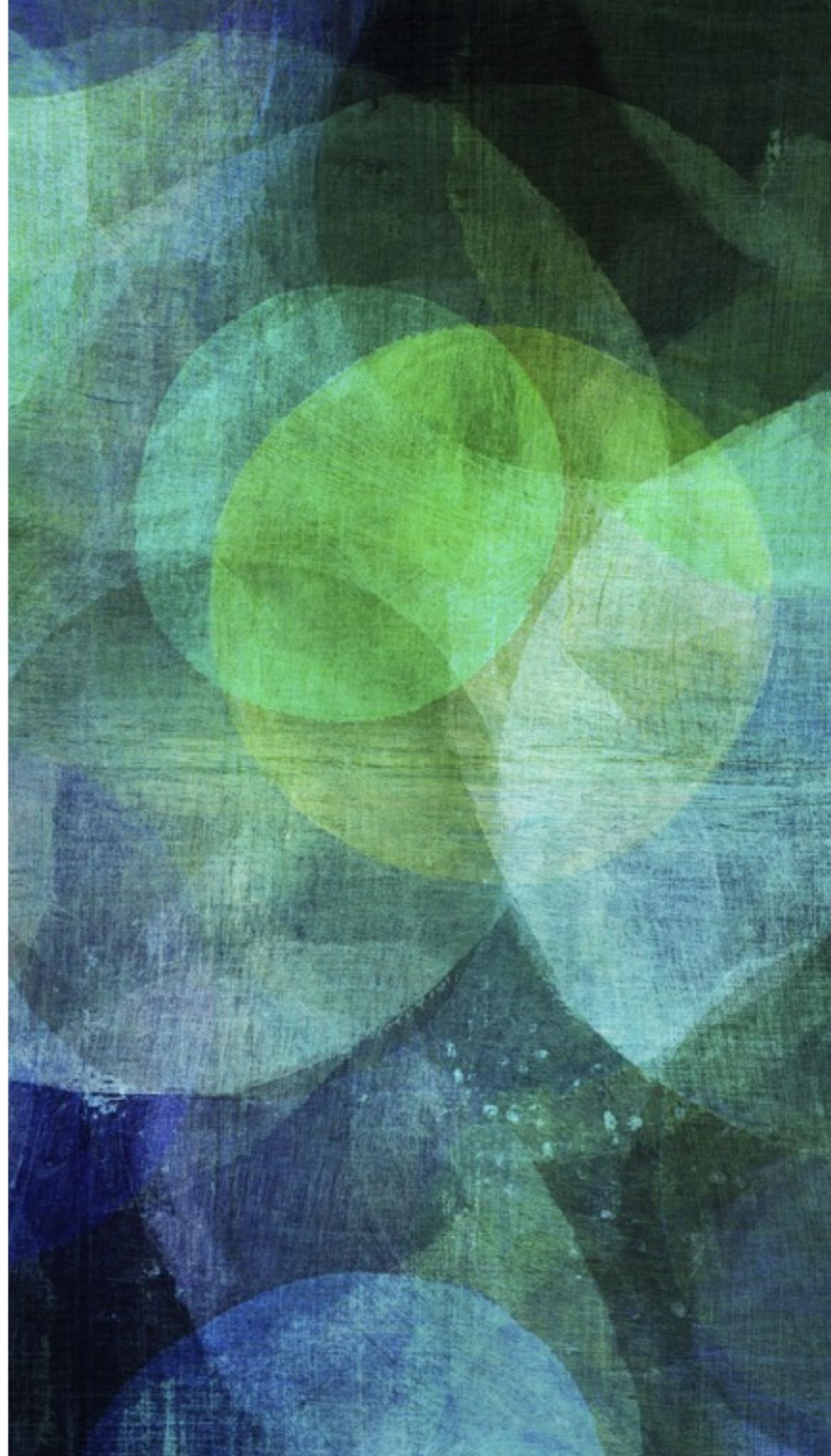
- `xprt_rdma.h`
 - `rpcrdma_req` – contains per-request state
 - `rpcrdma_rep` – state of a reply
 - `rpcrdma_mw` – state of one memory region
 - `rpcrdma_xprt` – per-transport state
 - `rpcrdma_regbuf` – internal buffer with DMA mapping state
 - `rpcrdma_buffer` – set of reqs and reps for one connection

LOCKING

- Implicit serialization
 - The RPC client serializes calls to `->send_request`, `->connect`, and `->close`
 - The provider serializes calls to completion handlers

Spin lock	Protected list
<code>rb_lock</code>	<code>rb_send_bufs</code> , <code>rb_recv_bufs</code> , <code>rb_pending</code>
<code>rb_mwlock</code>	<code>rb_mws</code> , <code>rb_all</code>
<code>rb_reqslock</code>	<code>rb_allreqs</code>
<code>rb_recoverylock</code>	<code>rb_stale_mrs</code>

SERVER TRANSPORT OVERVIEW



SERVER TRANSPORT SWITCH

- `svc_rdma_recvfrom.c`
 - `svc_rdma_recvfrom` – called by an `nfdsd` thread to receive an RPC message from a client and assemble it into an `xdr_buf`. Dequeues complete Receives, initiates RDMA Reads, dequeues complete Reads.
- `svc_rdma_sendto.c`
 - `svc_rdma_sendto` – called by an `nfdsd` thread to send an RPC message in an `xdr_buf` to a client. Initiates RDMA Writes and Sends.
- RDMA Read and Write WRs are scheduled in `svc_rdma_rw.c`

ACCEPTING CONNECTIONS

- `transport.c`
 - Sets up a listener QP
 - New connections accepted in `svc_rdma_accept`, which allocates fixed per-connection resources
- Some completion upcall handlers live in this file
- And one helper that posts Send operations

INTERESTING DATA STRUCTURES

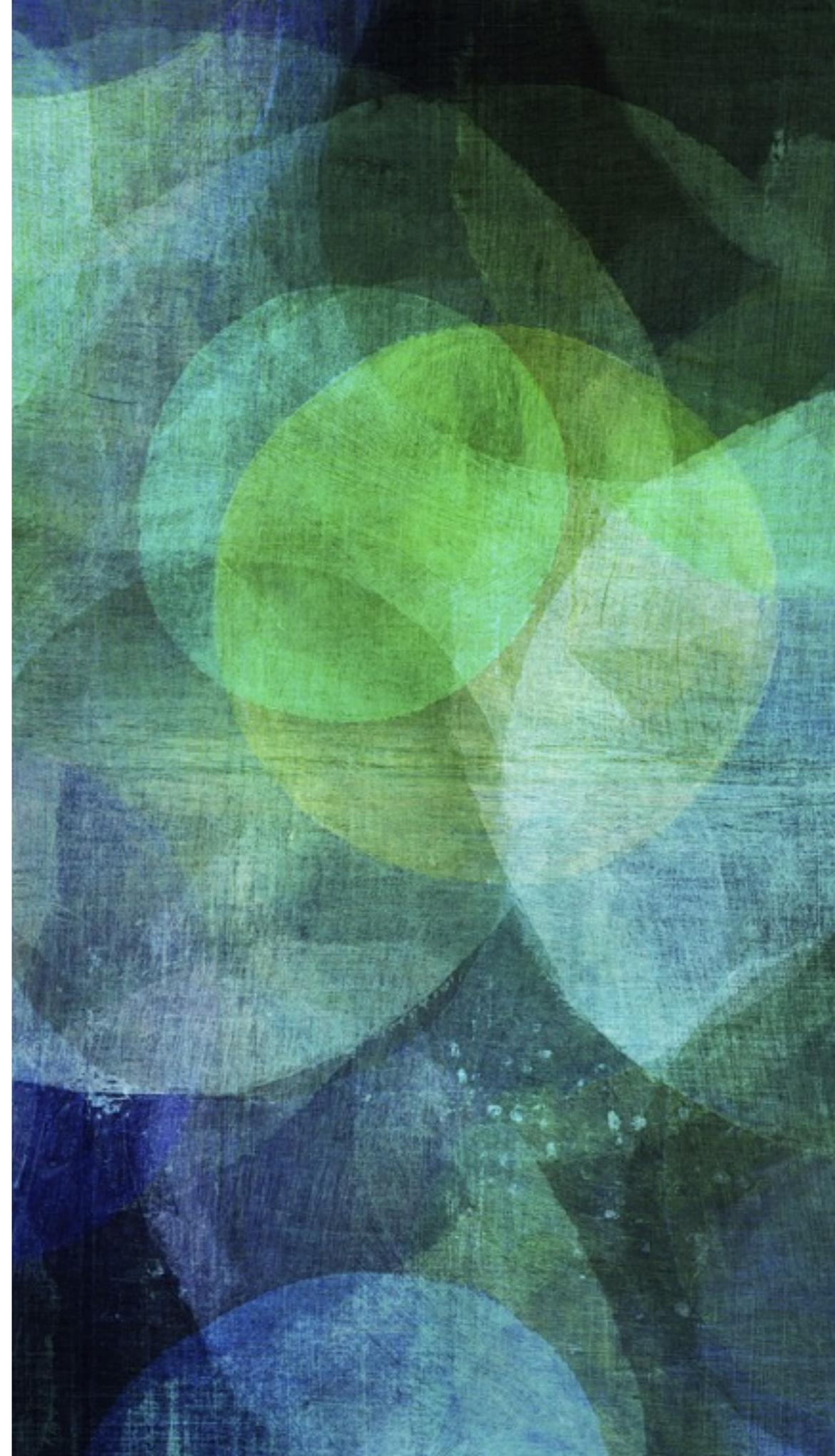
- `svcrdma_xprt` – per-connection state
- `svc_rdma_op_ctxt` – state for each Send and each Receive.
- `svc_rdma_rw_ctxt` – state for a set of RDMA Reads or Writes.
- `svc_rdma_chunk_ctxt` – completion context for one Read or Write chunk.
- `svc_rdma_write_info` – state for one Write chunk
- `svc_rdma_read_info` – state for one Read chunk

LOCKING

- Upper Layer locking
 - A per-transport mutex serializes calls to `->sendto`
 - Everything runs in a kthread or workqueue except `handle_connect_req`

Spin lock	Protected list
<code>sc_rq_dto_lock</code>	<code>sc_read_complete_q</code> , <code>sc_rd_dto_q</code>
<code>sc_ctxt_lock</code>	<code>sc_ctxts</code>
<code>sc_rw_ctxt_lock</code>	<code>sc_rw_ctxts</code>
<code>sc_lock</code>	<code>sc_accept_q</code>

NFSV4.1 BACK CHANNEL OPERATION



NFSV4.1 BACKCHANNEL

- `svc_rdma_backchannel.c`
 - Plugs into client transport switch
 - Sends CB Calls from the server, handles CB Replies

- `backchannel.c`
 - Plugs into RPC server framework
 - Handles incoming CB Calls on the client, sends CB Replies

LOCKING

Spin lock	Protected list
bc_pa_lock	bc_pa_list
rb_reqslock	rb_allreqs

